

# Evolutionary Bargaining



Game theory: Two players alternative offering game  
 Subgame perfect equilibrium found



Slight game modification →  
 Laborious work on new solutions  
 Perfect rationality assumption

Technical details (non-trivial)  
 Co-evolution  
 Incentive methods invented

- ◆ Proposal: EC for approximating solutions on new games

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# Two players alternative offering game

Player 1: How about 70% for me 30% for you?

$t = 0$ , Player 1's pay off is 70%

Player 1: No, how about 50-50?

$t = 2$ , Player 1's pay off is  $50\% \times e^{-0.1 \times 2} = 41\%$

- ◆ If neither players have any incentive to compromise, this can go on for ever
- ◆ Payoff drops over time – incentive to compromise
- ◆ A's Payoff =  $x_A \exp(-r_A t \Delta)$      Let  $r_1$  be 0.1,  $\Delta$  be 1

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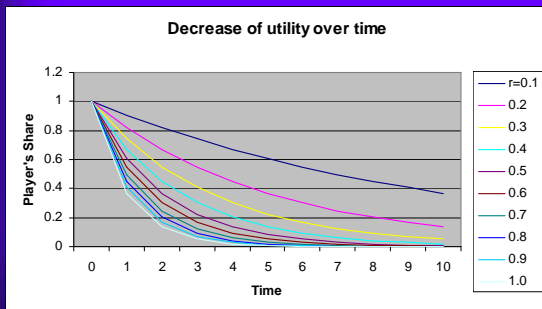
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# Payoff decreases over time



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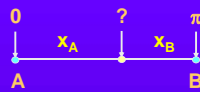
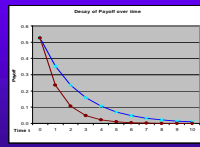
# Bargaining in Game Theory

- In reality: Offer at time  $t = f(r_A, r_B, t)$   
Is it necessary? (What is rational?)
- $t$  = # of rounds, at time  $\Delta$  per round
- A's payoff  $x_A$  drops as time goes by  
A's Payoff =  $x_A \exp(-r_A t\Delta)$
- Important Assumptions:
  - Both players rational
  - Both players know *everything*
- Equilibrium solution for A:  

$$\mu_A = (1 - \delta_B) / (1 - \delta_A \delta_B)$$
 where  $\delta_i = \exp(-r_i \Delta)$

Optimal offer:  $x_A = \mu_A$  at  $t=0$

Notice: No time  $t$  here

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# Evolutionary Computation for Bargaining

## Technical Details




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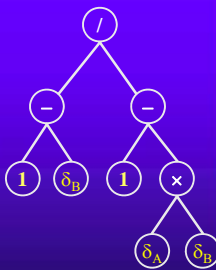
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# Issues Addressed in EC for Bargaining

- Representation
  - Should  $t$  be in the language?
- One or two population?
- How to evaluate fitness
  - Fixed or relative fitness?
- How to contain search space?
- Discourage irrational strategies:
  - Ask for  $x_A > 1$ ?
  - Ask for more over time?
  - Ask for more when  $\delta_A$  is low?



A run through

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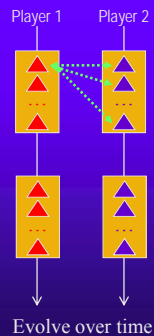
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## Two populations – co-evolution

- ◆ We want to deal with asymmetric games
  - E.g. two players may have different information
- ◆ One population for training each player's strategies
- ◆ Co-evolution, using relative fitness
  - Alternative: use absolute fitness



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## Representation of Strategies

- ◆ A tree represents a mathematical function  $g$
- ◆ Terminal set:  $\{1, \delta_A, \delta_B\}$
- ◆ Functional set:  $\{+, -, \times, \div\}$
- ◆ Given  $g$ , player with discount rate  $r$  plays at time  $t$ 
$$g \times (1 - r)^t$$
- ◆ Language can be enriched:
  - Could have included  $e$  or time  $t$  to terminal set
  - Could have included power  $^$  to function set
- ◆ Richer language  $\rightarrow$  larger search space  $\rightarrow$  harder search problem

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## Incentive Method: Constrained Fitness Function

- ◆ No magic in evolutionary computation
    - Larger search space  $\rightarrow$  less chance to succeed
  - ◆ Constraints are heuristics to focus a search
    - Focus on space where promising solutions may lie
  - ◆ Incentives for the following properties in the function returned:
    - The function returns a value in  $(0, 1)$
    - Everything else being equal, lower  $\delta_A \rightarrow$  smaller share
    - Everything else being equal, lower  $\delta_B \rightarrow$  larger share
- Note: this is the key to our search effectiveness

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## Incentives for Bargaining

◆  $F(g_i) =$

$\frac{GF(s(g_i)) + B}{B \geq 3 \text{ (tournament selection is used)}}$	If $g_i$ in $(0, 1]$ & $\underline{SM}_i > 0$ & $\underline{SM}_i > 0$
$\frac{GF(s(g_i)) + \underline{ATT}(i) + \underline{ATT}(i)}$	If $g_i$ in $(0, 1]$ & ( $\underline{SM}_i \leq 0$ or $\underline{SM}_i \leq 0$ )
$\frac{\underline{ATT}(i) + \underline{ATT}(i) - e^{(-1/ g_i )}}$	If $g_i$ NOT in $(0, 1]$

- ◆  $GF(s(g_i))$  is the game fitness (GF) of a strategy (s) based on the function  $g_i$ , which is generated by genetic programming
- ◆ When  $g_i$  is outside  $(0, 1]$ , the strategy does not enter the games

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## C1: Incentive for Feasible Solutions

- ◆ The function returns a value in  $(0, 1]$
- ◆ The function participates in game plays
- ◆ The game fitness (GF) is measured
- ◆ A bonus (B) incentive is added to GF
  - B is set to 3
  - Since tournament is used in selection, the absolute value of B does not matter (as long as  $B > 3$ )

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## C2: Incentive for Rational $\delta_A$

- ◆ Everything else being equal, lower  $\delta_A \rightarrow$  smaller share for A

Given a function  $g_i$ :

- ◆ The sensitive measure  $\underline{SM}_i(\delta_i, \delta_j, \alpha)$  measures how much  $g_i$  decreases when  $\delta_i$  increases by  $\alpha$

Attribute  $\underline{ATT}(i) =$

1	If $\underline{SM}_i(\delta_i, \delta_j, \alpha) > 1$
$-e^{(1/\underline{SM}_i(\delta_i, \delta_j, \alpha))}$	If $\underline{SM}_i(\delta_i, \delta_j, \alpha) \leq 1$

$0 < \underline{ATT}(i) \leq 1$

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### C3: Incentive for Rational $\delta_B$

- Everything else being equal, lower  $\delta_B \rightarrow$  larger share for A

Given a function  $g_i$ :

- The sensitive measure  $SM_i(\delta_i, \delta_j, \alpha)$  measures how much  $g_i$  increases when  $\delta_j$  increases by  $\alpha$

Attribute  $ATT(j) =$

1	If $SM_j(\delta_i, \delta_j, \alpha) > 1$
$-e^{(1/SM_j(\delta_i, \delta_j, \alpha))}$	If $SM_j(\delta_i, \delta_j, \alpha) \leq 1$

$$0 < ATT(j) \leq 1$$

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### Sensitivity Measure (SM) for $\delta_i$

- $SM_i(\delta_i, \delta_j, \alpha) =$

$\frac{[g_i(\delta_i \times (1 + \alpha), \delta_j) - g_i(\delta_i, \delta_j)]}{\div g_i(\delta_i, \delta_j)}$	If $\delta_i \times (1 + \alpha) < 1$
$\frac{[g_i(\delta_i, \delta_j) - g_i(\delta_i \times (1 - \alpha), \delta_j)]}{\div g_i(\delta_i, \delta_j)}$	If $\delta_i \times (1 + \alpha) \geq 1$

- $SM_i$  measures how much  $g_i$  decreases when  $\delta_i$  increases by  $\alpha$   
( $g_i$  is the function which fitness is to be measured)

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### Sensitivity Measure (SM) for $\delta_j$

- $SM_j(\delta_i, \delta_j, \alpha) =$

$\frac{[g_i(\delta_i, \delta_j) - g_i(\delta_i, \delta_j \times (1 + \alpha))]}{\div g_i(\delta_i, \delta_j)}$	If $\delta_j \times (1 + \alpha) < 1$
$\frac{[g_i(\delta_i, \delta_j \times (1 - \alpha)) - g_i(\delta_i, \delta_j)]}{\div g_i(\delta_i, \delta_j)}$	If $\delta_j \times (1 + \alpha) \geq 1$

- $SM_j$  measures how much  $g_i$  increases when  $\delta_j$  increases by  $\alpha$   
( $g_i$  is the function which fitness is to be measured)

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## Bargaining Models Tackled

Determinants	Complete Information	Uncertainty	
		1-sided	2-sided
Discount Factors	* Rubinstein 82	* Rubinstein 85 * <i>Imprecise info Ignorance</i>	* <i>Bilateral ignorance</i>
+ Outside Options	* Binmore 85	* <i>Uncertainty + Outside Options</i>	More could be done easily

\* = Game theoretical solutions known  
 x = game theoretic solutions unknown

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## Models with known equilibriums

### Complete Information

- ◆ **Rubinstein 82 model:**
  - Alternative offering, both A and B know  $\delta_A$  &  $\delta_B$
- ◆ **Binmore 85 model, outside options:**
  - As above, but each player has an outside offer,  $w_A$  and  $w_B$

### Incomplete Information

- ◆ **Rubinstein 85 model:**
  - B knows  $\delta_A$  &  $\delta_B$
  - A knows  $\delta_A$
  - A knows  $\delta_B$  is  $\delta_w$  with probability  $w_0$ ,  $\delta_s$  ( $> \delta_w$ ) otherwise

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## Models with unknown equilibriums

### Modified Rubinstein 85 / Binmore 85 models:

- ◆ **1-sided Imprecise information**
  - B knows  $\delta_A$  &  $\delta_B$ ; A knows  $\delta_A$  and a normal distribution of  $\delta_B$
- ◆ **1-sided Ignorance**
  - B knows both  $\delta_A$  and  $\delta_B$ ; A knows  $\delta_A$  but not  $\delta_B$
- ◆ **2-sided Ignorance**
  - B knows  $\delta_B$  but not  $\delta_A$ ; A knows  $\delta_A$  but not  $\delta_B$
- ◆ **Rubinstein 85 + 1-sided outside option**

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## Equilibrium with Outside Option

$x_A^*$	Conditions	
$\underline{\mu}_A$	$w_A \leq \delta_A \mu_A$	$w_B \leq \delta_B \mu_B$
$1 - w_B$	$w_A \leq \delta_A (1 - w_B)$	$w_B > \delta_B \mu_B$
$\delta_B w_A + (1 - \delta_B)$	$w_A > \delta_A \mu_A$	$w_B \leq \delta_B (1 - w_A)$
$1 - w_B$	$w_A > \delta_A (1 - w_B)$	$w_B > \delta_B (1 - w_A)$
$w_A$	$w_A + w_B > 1$	–

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## Equilibrium in Uncertainty – Rub85

$V_s = \frac{1 - \delta_s}{1 - \delta_1 \delta_s}$	$\delta_2 = \delta_w$		$\delta_2 = \delta_s$	
	$x_1^*$	$t^*$	$x_1^*$	$t^*$
$w_0 < w^*$	$V_s$	0	$V_s$	0
$w_0 > w^*$	$x^{w_0}$	0	$\frac{1 - ((1 - x^{w_0}) / \delta_w)}{1 - \delta_1 \delta_s}$	1

$$w^* = \frac{V_s - \delta_1^2 V_s}{1 - \delta_w + \delta_1 V_s (\delta_w - \delta_1)}$$

$$x^{w_0} = \frac{(1 - \delta_w)(1 - \delta_1^2(1 - w_0))}{1 - \delta_1^2(1 - w_0) - \delta_1 \delta_w w_0}$$

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## Evolutionary Bargaining Conclusions

- ◆ **Demonstrated GP's flexibility**
  - Models with known and unknown solutions
  - Outside option
  - Incomplete, asymmetric and limited information
- ◆ **Co-evolution is an *alternative approximation method* to find game theoretical solutions**
  - Perfect rationality assumption relaxed
  - Relatively quick for approximate solutions
  - Relatively easy to modify for new models
- ◆ **Genetic Programming with incentive / constraints**
  - Constraints helped to focus the search in promising spaces
- ◆ **Lots remain to be done...**

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# Running GP in Bargaining

Representation, Evaluation  
Selection, Crossover, Mutation

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## Representation

◆ Given  $\delta_A$  and  $\delta_B$ , every tree represents a constant

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## Population Dynamics

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## Evaluation

- ◆ Given the discount factors, each tree is translated into a constant  $x$ 
  - It represents the demand represented by the tree.
- ◆ All trees where  $x < 0$  or  $x > 1$  are evaluated using rules defined by the incentive method
- ◆ All trees where  $0 \leq x \leq 1$  enter game playing
- ◆ Every tree for Player 1 is played against every tree for Player 2

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## Evaluation Through Bargaining

Demands by Player 2's strategies

		.46	.31	.65	.20	Player 1 Fitness
Player 1 Demands	.75	0	0	0	.75	0.75
	.24	.24	.24	.24	.24	0.96
	.36	.36	.36	0	.36	1.08
	.59	0	.59	0	.59	1.18

- ◆ Incentive method ignored here for simplicity

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## Selection

Rule (Demand)	Fitness	Normalized	Accumulated
R1 (0.75)	0.75	0.19	0.19
R2 (0.96)	0.96	0.24	0.43
R3 (1.08)	1.08	0.27	0.70
R4 (1.18)	1.18	0.30	1.00
Sum:	3.97	1	

- ◆ A random number  $r$  between 0 and 1 is generated
- ◆ If, say,  $r=0.48$  (which is  $>0.43$ ), then rule R3 is selected

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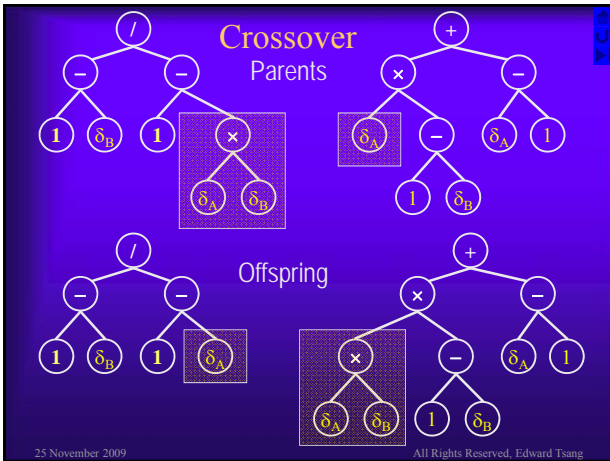
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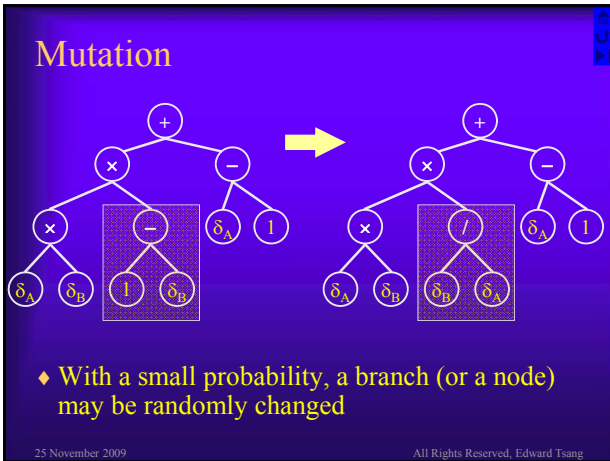
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◆ With a small probability, a branch (or a node) may be randomly changed

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